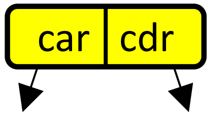
Functional Data Structures

[C. Okasaki, *Simple and efficient purely functional queues and deques*, J. of Functional Programming, 5(4), 583-592, 1995]
[H. Kaplan, R. Tarjan, *Purely functional, real-time deques with catenation*, Journal of the ACM, 46(5), 577-603, 1999]

Purely functional



never modify
only create new pairs
only DAGs

(Atomic values: Integers, Chars, Float, Bool, ...)

Strict evaluation
Evaluate list now

Lazy evaluation/memoization
First add element when head needed and return function *lazy* incrementing the rest

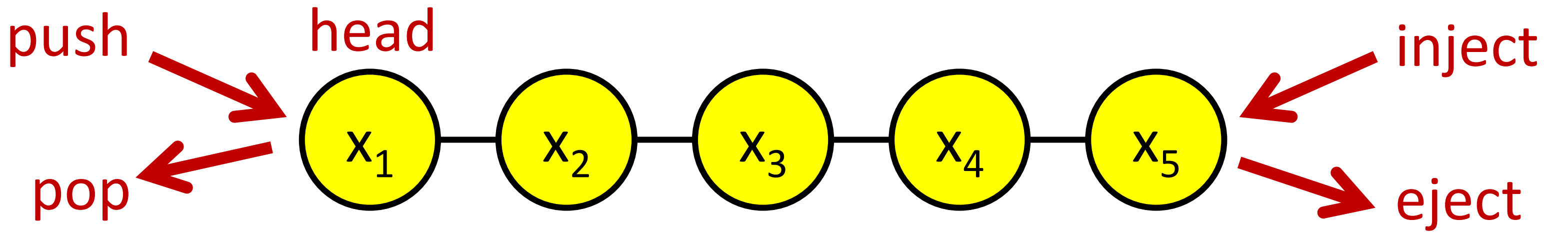
Example

```
inc(())     = ()
inc(e::L')  = (e+1)::inc(L')
```

List operations

Deque
= Double Endede Queue
(Donald E. Knuth 74)

	Stack	Queue	Deque	Catenable lists	Catenable deques
makelist(x)	★	★	★	★	★
push(x,L)	★		★	★	★
pop(L)	★	★	★	★	★
inject(x,L)		★	★		★
eject(L)			★		★
catenate(K,L)				★	★



Catenable lists (slow)

```
cat((),L)      = L
cat(e::K,L)    = e::cat(K,L)
```

} O(length 1st list)

List reversal

```
rev(L)         = rev′(L,())
rev′((),T)     = T
rev′(e::L,T)   = rev′(L,e::T)
```

} O(|L|)

Bad if expensive operation repeated

Queues (Head,Tail) **Ex: ((1,2,3),(5,4)) ≡ [1,2,3,4,5]**

```
inject(e,(H,T)) = (H,e::T)
```

} O(1)

Version 1

```
pop((e::H,T))  = (e,(H,T))
pop(((),T))    = (e,(T′,())) where e::T′ = rev(T)
```

} Strict
O(1) amortized
$\Phi = |T|$

Version 2 (Invariant |H|≥|T|)

```
pop((e::H,T))  = (e,(H,T))                   if |H|≥|T|
               = (e,(cat(H,rev(T)),()))if |H|<|T|
Inject(e,(H,T)) = (H,e::T)                    if |H|>|T|
                = (cat(H,rev(e::T)),())if |H|≤|T|
```

} Lazy **Good**
O(1) amortized

```
cat((),L)        = L
cat(e::K,L)      = e::cat(K,L)
```

lazy evaluation $\rightarrow$ recursive call first evaluated when 1st element accessed

🙂

```
rev(L)           = rev'(L,())
rev'((),T)       = T
rev'(e::L,T)     = rev'(L,e::T)
```

lazy evaluation $\rightarrow$ everything evaluated when 1st element accessed

🙁

```
inject(e,(H,T)) = (H,e::T)
```

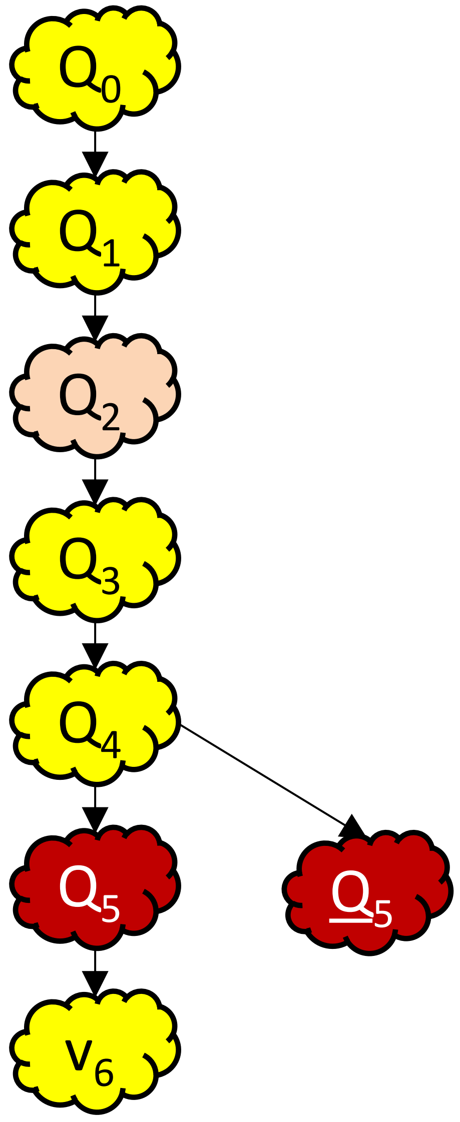
TRICK In `cat(H,rev(T))` the cost for rev (T) is paid by the subsequent pops (with no reversals) from the H part of the catenation. All pops deleting from H pays O(1) for doing O(1) work of the reverse.

lazy evaluation

Version 2 (Invariant |H|≥|T|)

```
pop((e::H,T)) = (e,(H,T))                 if |H|>|T|
              = (e,(cat(H,rev(T)),()))    if |H|≤|T|
```

[C. Okasaki, *Simple and efficient purely functional queues and deques*, J. of Functional Programming, 5(4), 583-592, 1995]



setup expensive
lazy evaluation

cheap operations
amortize cost of upcomming
expensive lazy evaluation

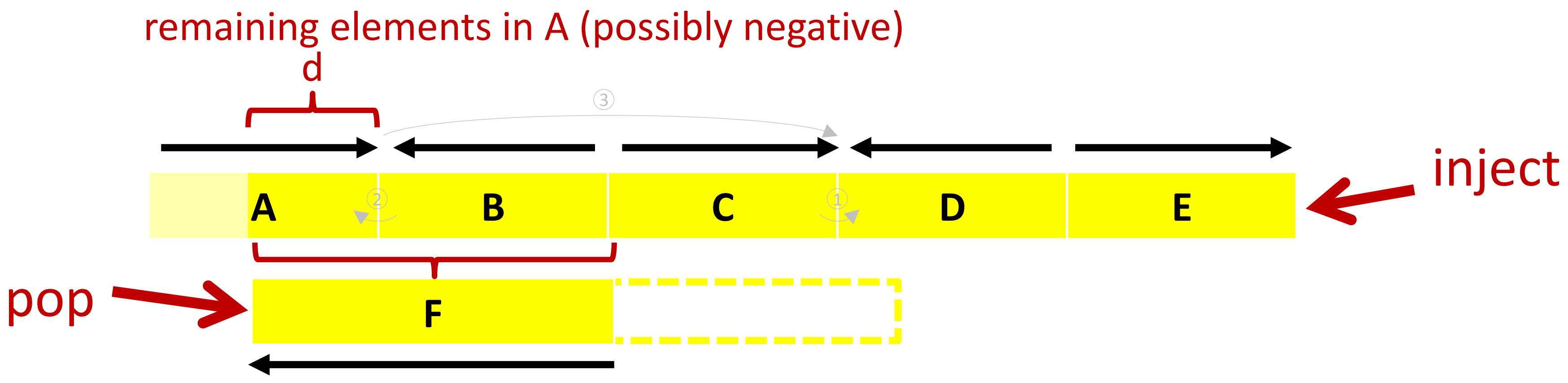
execute
expensive

only **one** of
Q_5 or $\underline{Q}_5$
is expensive

Real-time Queues i.e. strict worst-case O(1) time

[R. Hood, R. Melville, *Real-time queue operations in pure Lisp*. Information Processing Letters, 13, 50-54, 1981]

- incremental version of the amortized solution



```
                                d  F   A   B   C   D  E
makelist(x)                  = (0,(x),(),(x),(),(),())
inject(x,(d,F,A,B,C,D,E))    = f(f(d,F,A,B,C,D,x::E))
pop((0,F,A,(),(),x::D,E))    = (x,f(0,D,(),D,E,(),()))
pop((d,x::F,A,B,C,D,E))      = (x,f(f(d-1,F,A,B,C,D,E)))

① f(d,F,A,B,x::C,D,E)     = (d,F,A,B,C,x::D,E)
② f(d,F,A,x::B,C,D,E)     = (d+1,F,x::A,B,C,D,E)
③ f(d,F,x::A,(),(),D,E)   = (d-1,F,A,(),(),x::D,E)  if d>0
  f(d,F,A,(),(),D,E)      = (0,D,(),D,E,(),())      if d=0
```

Queue = ABCDE with first |A|-d removed

F = prefix of queue with $|F| \geq d+|B|$

$0 \leq d+(|B|-|C|)/2$

$|E|+|A|/2 \leq |D|+d$

Queues

[R. Hood, R. Melville, *Real-time queue operations in pure Lisp*. Information Processing Letters, 13, 50-54, 1981] Strict, worst-case O(1)

[C. Okasaski, *Simple and efficient purely functional queues and deques*. Journal of Functional Programming 5,4, 583-592, 1995] Lazy, amortized O(1)

Catenable lists

[S.R. Kosaraju, *Real-time simulation of concatenable double-ended queues by double-ended queues*, Proc. 11th Annual ACM Symposium on Theory of Computing, 346-351, 1979]

[S.R. Kosaraju, *An optimal RAM implementation of catenable min double-ended queues*, Proc. 5th Annual ACM-SIAM Symposium on Discrete Algorithms, 195-203, 1994]

} Not confluently persistent

[J.R. Driscoll , D.D. Sleator , R.E. Tarjan, *Fully persistent lists with catenation*, Journal of the ACM, 41(5), 943-959, 1994] O(loglog k)

[A.L. Buchsbaum , R.E. Tarjan, *Confluently persistent deques via data-structural bootstrapping*, Journal of Algorithms, 18(3), 513-547, 1995] { $2^{O(\log^* k)}$ / O(log* k)

} Not funtional

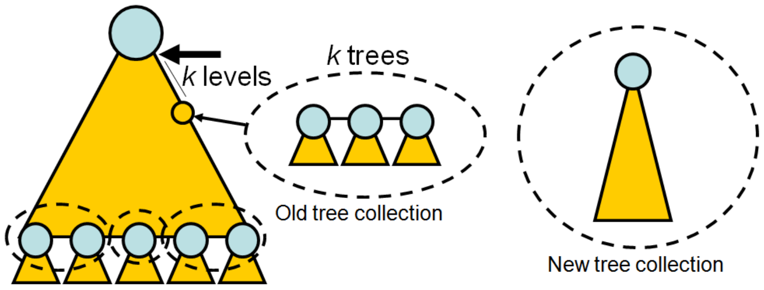
[H. Kaplan, R. Tarjan, *Purely functional, real-time deques with catenation*, Journal of the ACM, 46(5), 577-603, 1999] Strict, worst-case O(1)

[H. Kaplan, C. Okasaki, R.E. Tarjan, *Simple Confluently Persistent Catenable Lists*, SIAM Journal of Computing 30(3), 965-977 (2000)] Lazy, amortized O(1)

Functional Concatenable Search Trees

[G.S. Brodal, C.Makris, K. Tsichlas, *Purely Functional Worst Case Constant Time Catenable Sorted Lists*, In Proc. 14th Annual European Symposium on Algorithms, LNCS 4168, 172-183, 2006]

- Search, update O(log n)
- Catenation O(1)



Open problems

- Split O(log n) ?
- Finger search trees with O(1) time catenation ?
- Search trees with O(1) space per update ?